



Contextual Search in the Presence of Irrational Agents

Akshay Krishnamurthy (MSR NYC), Thodoris Lykouris (MSR NYC → MIT), Chara Podimata (Harvard), and Robert Schapire (MSR NYC)

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Full version:



Main Question

Contextual search $\hat{=}$ generalization of single dimensional binary search where a context x_t is given per round to guide the search

Can model many real-life settings, e.g.,
AirBnB's smart pricing algorithm



How do we construct contextual search algorithms that are robust to the presence of some **model misspecifications**?

TCS: multidimensional binary search when nature is allowed to lie a fixed number of times

ML: learning with 1-sided feedback in the presence of corruptions

Econ: feature-based pricing against adversarially irrational agents (i.e., deviating from prescribed behavioral model)

Model [Cohen, Lobel, Paes Leme, EC16]

Total corruption (unknown to learner): $C = \sum_t c_t$.

Nature selects hidden parameter $\theta^* \in \mathbb{R}^d$.

For round $t \in [T]$:

1. Nature chooses observable context $x_t \in \mathbb{R}^d$ with $\|x_t\|_2 \leq 1$.

and corruption $c_t \in \{0, 1\}$

2. Learner selects a price $p_t \geq 0$.

3. Agent makes purchase if $\langle x_t, \theta^* \rangle \geq p_t$ if $c_t = 0$

or arbitrarily if $c_t = 1$

4. Learner observes if purchase occurred (*one-sided feedback*).

5. Learner collects revenue $p_t \cdot \mathbb{1}\{\langle x_t, \theta^* \rangle \geq p_t\}$.

 $Regret = \sum_t \langle x_t, \theta^* \rangle - \sum_t p_t \cdot \mathbb{1}\{\langle x_t, \theta^* \rangle \geq p_t\}$

[Lobel, Paes Leme, Vladu, EC17]: $Regret = O(d \log T)$

Main Result

Robust contextual search algorithm **agnostic to C** achieving $Regret = O(Cd^3 \log^3 T)$ (w. high prob.)

[Ulam 76], [Rivest, Meyer, Kleitman, Winkmann, Spencer, JCSS80], [Karp and Kleinberg, SODA07]
For Ulam's game ($d = 1$) tight bound of $\log T + C \log C + C \log \log T$ [RMKWS, JCSS80]

Extends classical Ulam's game (aka twenty questions with a liar) in 2 ways: 1. Multidimensional 2. Agnostic to C

Complications of Adversarial Irrationality

Binary search in 1d: after at most $\log(\frac{1}{\epsilon})$ queries identify the hidden parameter θ^* within an ϵ interval.



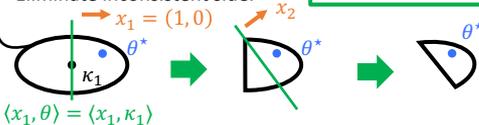
What if one of the adversary's responses is corrupted?

Completely wrong interval returned after $\log(\frac{1}{\epsilon})$
A way out: Repeat queries!

Non-Robust Contextual Search Algorithm

[Lobel, Paes Leme, Vladu]: $Regret = O(d \log T)$

- Maintain active knowledge set with feasible values for θ^* .
- Set price $\langle \kappa_t, x_t \rangle$ (κ_t : centroid).
- Eliminate inconsistent side.



Important properties of cut

- 1. Never eliminate θ^* .
- 2. Volumetric progress

Challenges and Solutions

1. Cannot repeat same query (adversarially chosen x_t 's)

→ Keep "penalty" for each θ . Make cut once region with penalty $\leq C$ (protected region, $\mathcal{P}(C)$) is fully on one side

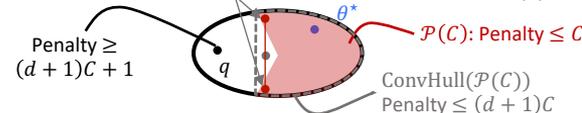
2. May never have a context cut with $\mathcal{P}(C)$ fully on one side (even for $C = 1$ and infinitely many contexts).

→ Run algorithm in epochs.

→ Combine context cuts to a valid cut (w. properties in green box).

→ Epoch size to guarantee existence of valid cut: $2d(d+1)C + 1$.

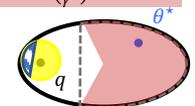
[Caratheodory's theorem]: every point $p \in \text{ConvHull}(\mathcal{P}(C))$ can be written as the convex combination of $\leq d + 1$ points in $\mathcal{P}(C)$.



Given q and $\text{ConvHull}(\mathcal{P}(C)) \rightarrow$ Perceptron finds valid cut.

3. Perceptron needs a margin γ to converge after $O(\frac{1}{\gamma^2})$ mistakes.

- Sample random point q from ball (enough penalty for its center).
- W. const. prob., q has enough margin.



→ Use Perceptron as a witness:

- If q has enough margin, Perceptron terminates after $O(\frac{1}{\gamma^2})$ rounds.
- Else, resample

4. From known to unknown C.

Subsampling and global eliminations of $\log T$ layers.

→ Extension of the multi-layering race technique of [Lykouris, Mirrokni, Paes Leme, STOC18] for continuous action spaces.

Known C