

# Learning for Dynamic Pricing

# Online Posted-Price Mechanisms

Assume the mechanism designer (MD) has infinite copies of the same item.

At round  $t = 1, \dots, T$ :

1. MD selects price  $p_t$ .
2. An agent with (hidden) valuation  $v_t$  arrives.
3. Agent buys iff  $v_t \geq p_t$ .
4. MD observes  $1\{v_t \geq p_t\}$ .
5. MD gets revenue  $p_t \cdot 1\{v_t \geq p_t\}$ .

If agents come from a fixed valuation distribution  $D$  (i.e.,  $v_t \sim D$ ) then the price maximizing the mechanism designer's expected revenue is:

$$p^* = \max_{p \in [0,1]} p \cdot \Pr_D[v \geq p]$$

“demand curve”

? How should the MD/learner choose prices  $\{p_t\}_{t \in [T]}$  if demand curve is unknown?

[Kleinberg and Leighton, FOCS03]

[Kleinberg and Leighton, FOCS03]

Types of valuation sequences:

1. **Stochastic** from unknown distribution  $v_t \sim D, \forall t$ .
2. **Adversarial**:  $\{v_t\}_{t \in [T]}$  is chosen arbitrarily.
3. **Identical** (but unknown).

Goal: (Tight) Regret Minimization

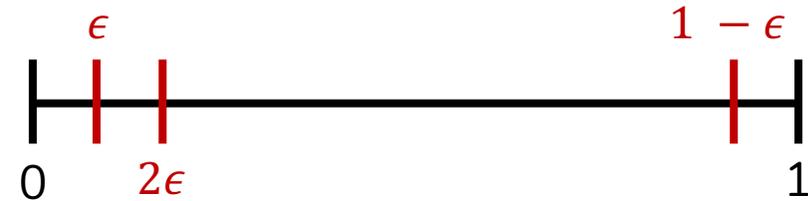
# Learning in Dynamic Pricing

Sequence of  $v_t$  chosen stochastically/adversarially.

At round  $t = 1, \dots, T$ :

1. Learner selects price  $p_t$ .
2. An agent with (hidden) valuation  $v_t$  arrives.
3. Agent buys iff  $v_t \geq p_t$ .
4. Learner observes  $1\{v_t \geq p_t\}$ .
5. Learner gets revenue  $p_t \cdot 1\{v_t \geq p_t\}$ .

Regret upper bound: uniform discretization



Nearly tight in worst case!

For adversarial  $\{v_t\}_{t \in [T]}$ : [[Podimata](#) and Slivkins, COLT21] proposed *adaptive discretization algorithm* that works despite the learner's revenue function not being (fully) Lipschitz.

[Kleinberg and Leighton, FOCS03]

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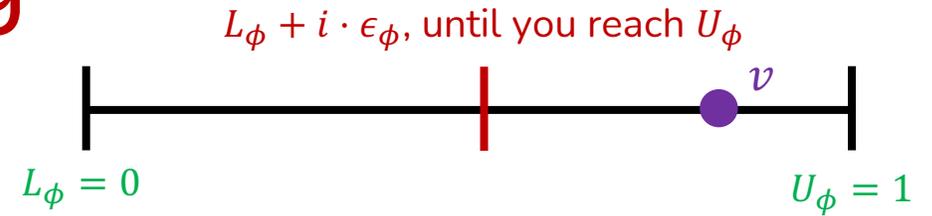
# Learning in Dynamic Pricing

Common  $v_t = v, \forall t \in [T]$ .

At round  $t = 1, \dots, T$ :

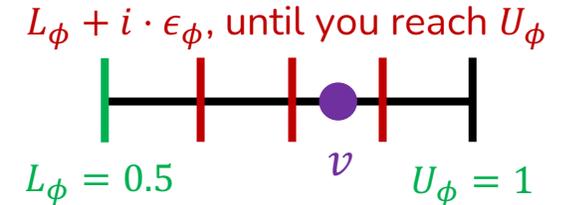
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Phase  $\phi = 1$   
 $\epsilon_\phi = 1/2$



If one of the prices was rejected:

Phase  $\phi = 2$   
 $\epsilon_\phi = \epsilon_{\phi-1}^2$



- Continue with phases until interval with length  $1/T$ .
- Then, play lowest price in interval forever.

$$R(T) = \Theta(\log \log T)$$

[Kleinberg and Leighton, FOCS03]

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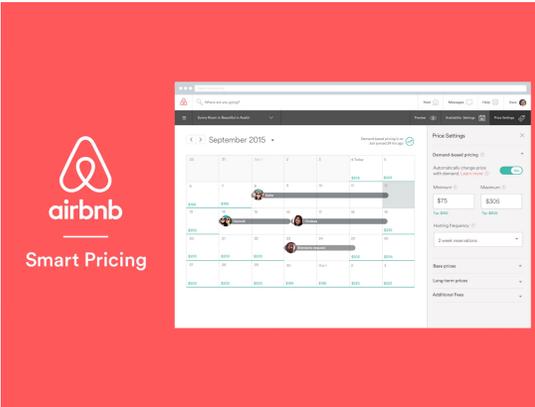
Goal: (Tight) Regret Minimization

# Contextual/Feature-Based Pricing

Repeat:

Nature chooses  $d$ -dimensional context  $x_t$

- 1. Product arrives with *features (location, #rooms, etc)*
- 2. Airbnb suggests *a price*    Learner chooses price  $p_t \in \mathbb{R}$ 
  - Exists common feature value  $\theta^* \in \mathbb{R}^d$ .
  - Agent **buys only iff** context valuation above price:  $\langle x_t, \theta^* \rangle \geq p_t$
- 3. User *purchases or not*
- 4. Airbnb *observes feedback (if purchase occurred)*    Learner's feedback:  $1\{\langle x_t, \theta^* \rangle \geq p_t\}$
- 5. Airbnb *collects revenue*    Learner's reward:  $p_t \cdot 1\{\langle x_t, \theta^* \rangle \geq p_t\}$



Set *prices* to collect as much *revenue* as possible.

*Regret* = Total revenue for benchmark pricing policy – Total revenue collected

# (Mostly) Related Work

Dynamic Pricing is a very important problem in OR; there are a ton of very interesting papers that I'm not listing here

- [Kleinberg and Leighton, FOCS03]: single dim, no context
- [Paes Leme, Sivan, Teng, Worah, ICML21]: single dim,  $v_t$ 's change by a little from round to round
- [Amin, Rostamizadeh, Syed NeurIPS14]: multidimensional, iid contexts
- [Cohen, Lobel, Paes Leme EC16/MS19], [Lobel, Paes Leme, Vladu EC17/OR18], [Paes Leme and Schneider, FOCS18], [Liu, Paes Leme, Schneider SODA21]: adversarial contexts, refined bounds
- [Mao, Paes Leme, Schneider NeurIPS18]: adversarial contexts, Lipschitz buyers
- [Krishnamurthy, Lykouris, Podimata, Schapire, STOC21]: adversarial contexts, some adversarially irrational buyers

# Behavioral Model Core Assumption

## Full Rationality

- Exists common feature value  $\theta^* \in \mathbb{R}^d$ .
- Agent **buys only iff** their context valuation above price:  $\langle x_t, \theta^* \rangle \geq p_t$

! But *sometimes* agents deviate from the prescribed behavioral model (**irrationality**)  
e.g., agents **may not buy**, even if price is below their valuation, **due to their psychological state**

## Problem

Model misspecifications can cause arbitrary failures to standard algorithms

## Main Question

How do we design **contextual search** algorithms that are **robust** to the presence of some **irrational** agents?

# Contextual Pricing for Fully Rational Agents

For round  $t = 1 \dots T$ :

1. Nature chooses observable context  $x_t \in \mathbb{R}^d$  with  $\|x_t\|_2 \leq 1$

[Cohen, Lobel, Paes Leme, EC16/MS19]

2. Learner selects a price  $p_t$

[Lobel, Paes Leme, Vladu, EC17/OR18]

3. Agent makes purchase if  $\langle x_t, \theta^* \rangle \geq p_t$

[Paes Leme, Schneider, FOCS18]

[Liu, Paes Leme, Schneider, SODA21]

4. Learner observes if purchase occurred (*one-sided feedback*)

5. Learner collects revenue if purchase occurred

$$\text{Regret} = O(d \log T)$$

$$\text{Benchmark pricing policy: } p_t^* = \langle x_t, \theta^* \rangle$$

- For revenue objective, [Liu, Paes Leme, Schneider, SODA21] improved to:  $O(d \log \log T)$
- For  $d = 1$ , optimal solution due to [Kleinberg and Leighton, FOCS03]

# Contextual Pricing for Adversarially Irrational Agents

Total corruption:  $C = \sum_t c_t$

For round  $t = 1 \dots T$ :

1. Nature chooses observable context  $x_t \in \mathbb{R}^d$  with  $\|x_t\|_2 \leq 1$

**and corruption  $c_t \in \{0, 1\}$**

2. Learner selects *a price*  $p_t$

3. Agent makes purchase if  $\langle x_t, \theta^* \rangle \geq p_t$  **if  $c_t = 0$**   
**or arbitrarily if  $c_t = 1$**

4. Learner observes if purchase occurred (*one-sided feedback*)

5. Learner **collects revenue** if purchase occurred

Extends the corruptions model studied in:  
[Lykouris, Mirrokni, Paes Leme STOC18]  
[Gupta, Koren, Talwar COLT19]  
[Zimmert, Seldin JMLR21]

## Desiderata for our Algorithms

- ✓ Graceful degradation of regret with  $C$
- ✓ No knowledge of  $C$  assumed (*agnostic*)

Total corruption:  $C = \sum_t c_t$

Number of features:  $d$

Robust contextual search algorithm **agnostic** to  $C$  achieving

$$\text{Regret} \leq C \cdot d^3 \cdot (\log T)^3 \text{ (with high probability)}$$

[Ulam 76], [Rivest, Meyer, Kleitman, Winklmann, Spencer, JCSS80], [Karp and Kleinberg, SODA07]  
For Ulam's game ( $d = 1$ ) tight bound of  $\log T + C \log C + C \log \log T$  [RMKWS, JCSS80]

Extends classical Ulam's game (aka twenty questions with a liar) in 2 ways:

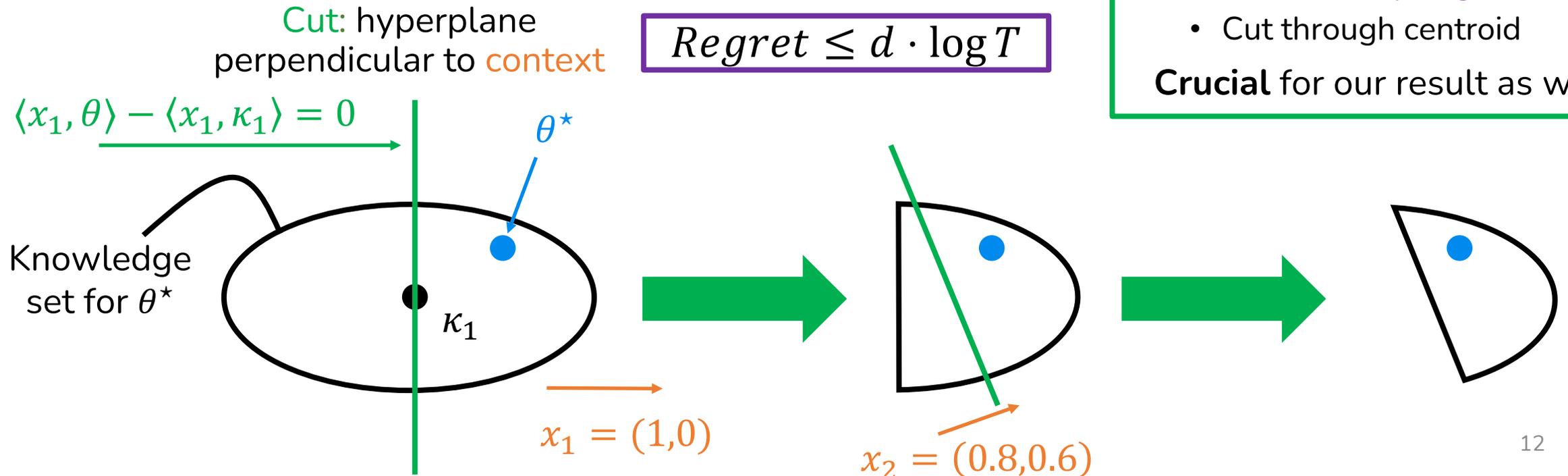
1. Multidimensional
2. Agnostic to  $C$

# The Non-Robust Algorithm

[Lobel, Paes Leme, Vladu, EC17/OR18]

## Multidimensional Binary Search

- Maintain active knowledge set with feasible values for  $\theta^*$ .
- Set price  $\langle \kappa_1, x_1 \rangle$  where  $\kappa_1$  centroid of knowledge set.
- Eliminate inconsistent side of knowledge set.



### Important properties of cut

1. Never eliminate  $\theta^*$ 
    - Retain all parameters consistent with feedback
  2. Volumetric progress
    - Cut through centroid
- Crucial** for our result as well

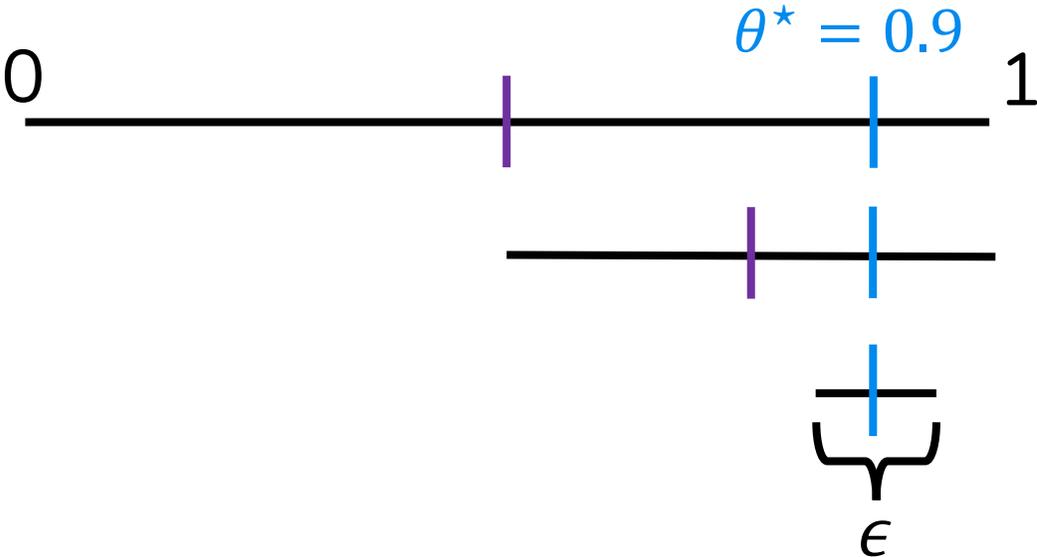
# Issue with Adversarial Irrationality: Ulam's Game

## Challenge

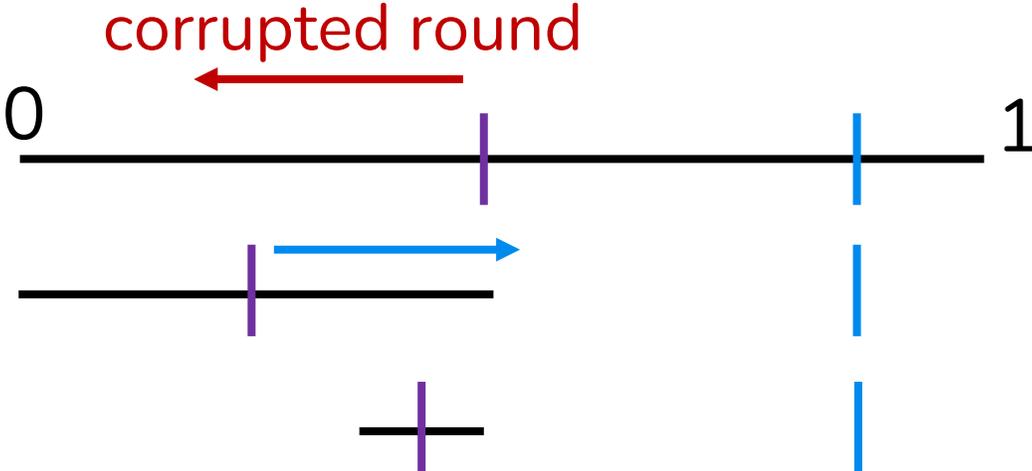
How can you fix it (in high dimensions) when you know that  $\bar{c}$  responses are corrupted?

### Binary search ( $d = 1$ )

- I fix a number  $\theta^* \in [0,1]$ .
- How many queries do you need to find it with accuracy  $\epsilon$ ?



... when *one response* is corrupted?

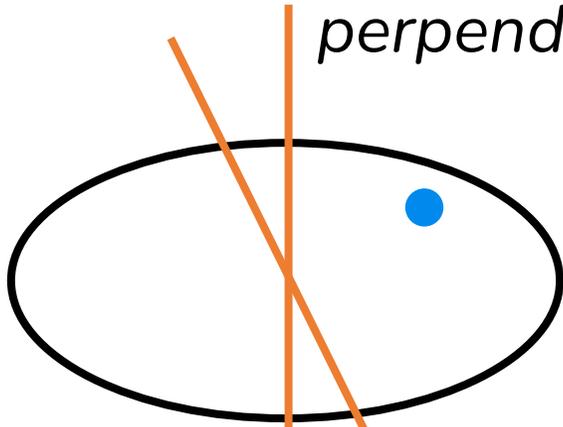


# Robust Volumetric Progress

## Challenge 1

We cannot repeat the same query (contexts are different at different rounds).

*Context cut: hyperplane perpendicular to context*



$x_1 = (1,0)$

$x_2 = (0.8,0.6)$

## Idea 1

Keep “penalty” for each parameter & make cut once a context cut fully retains protected region on one side

- Protected region: all parameters with “penalty”  $\leq \bar{c}$

# Robust Volumetric Progress

## Challenge 2

We may never have a **context cut** with the **protected region** fully on one side.

**Counterexample:** Even with infinite contexts and  $\bar{c} = 1$ , no such **context cut**.

# Robust Volumetric Progress

## Challenge 2

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## Idea 2

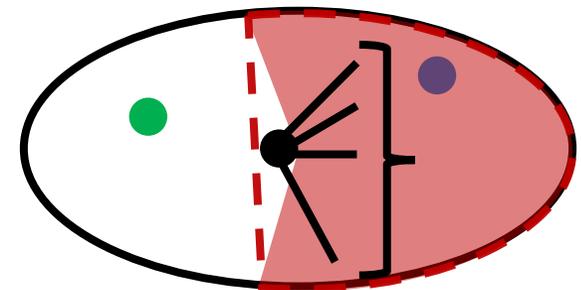
Combine **context cuts** to compute a “**valid cut**”.

## Important properties of valid cut

1. Never eliminate  $\theta^*$ 
  - Fully retains **protected region** on one side.
2. Volumetric progress
  - Cross **close** to centroid

## Idea 3

Show that  $2d \cdot (d + 1) \cdot \bar{c} + 1$  **context cuts** have enough information to compute **such a valid cut** (**Caratheodory's theorem**)



$d + 1$  points<sup>16</sup>

# Robust Volumetric Progress

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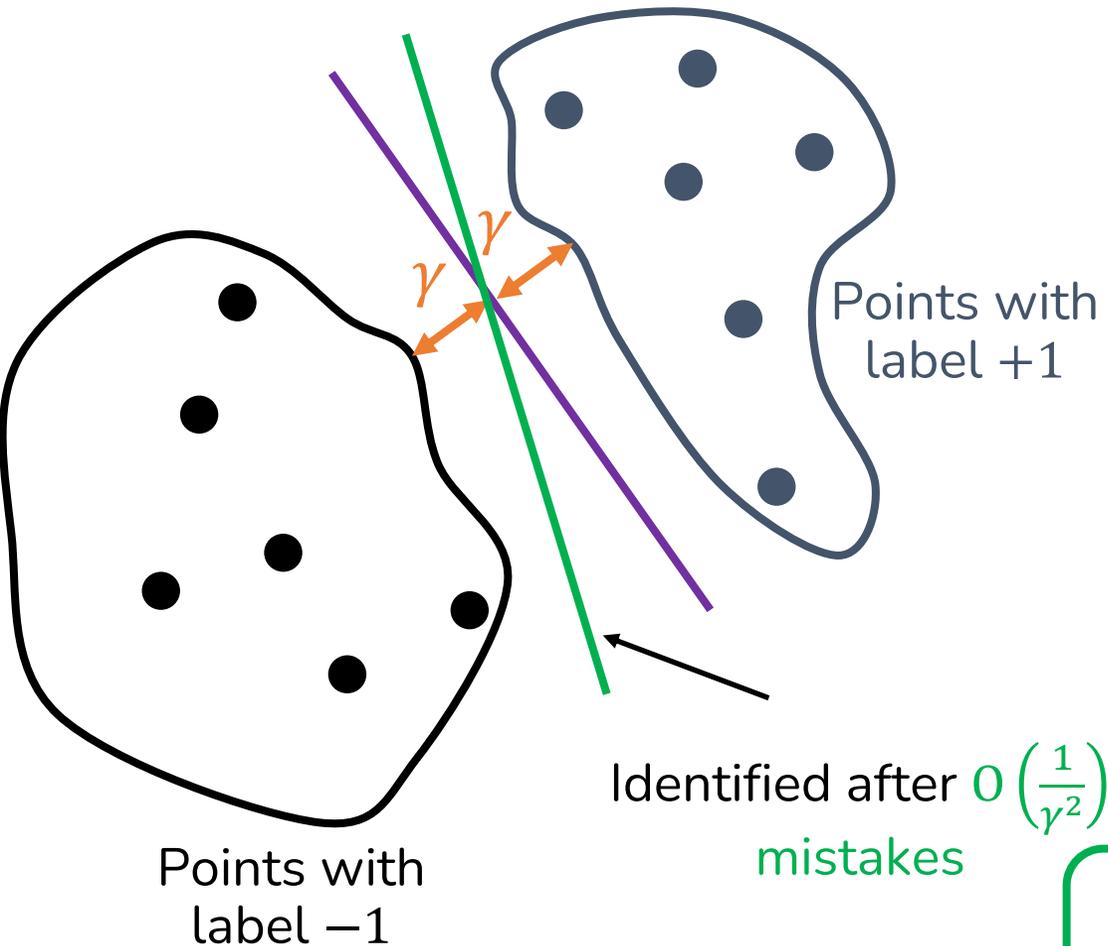
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## Idea 4

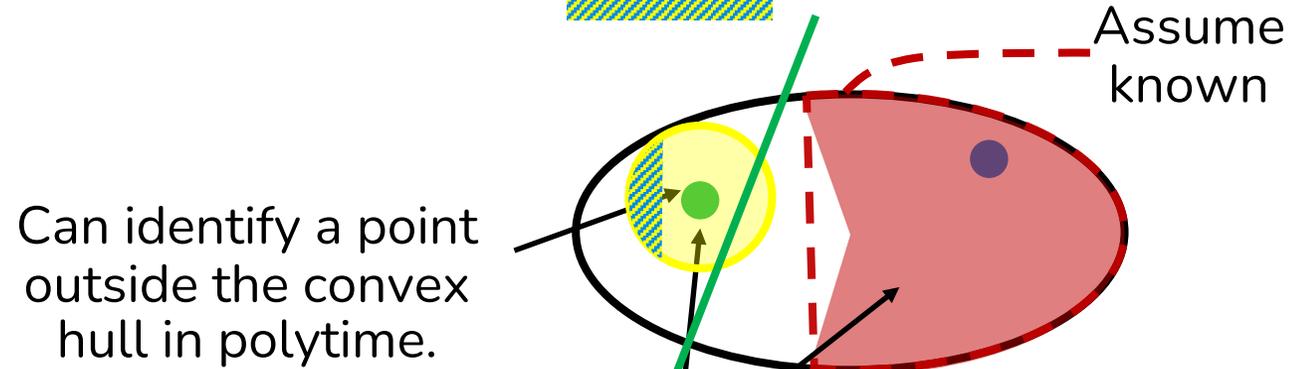
Use Perceptron as a **witness** for finding a **valid cut**.

# Perceptron as a Witness

## Standard Perceptron



- Sample randomly from ball
- With constant probability, you will sample point from pattern.



Idea

If we knew them exactly + they had margin  $\rightarrow$  Perceptron  $\rightarrow$  done

- If dataset correct  $\rightarrow$  Perceptron mistakes  $O\left(\frac{1}{\gamma^2}\right)$
- Else, resample

# Unknown $C$ : Multi-Layering Race to the Rescue

Original: [Lykouris, Mirrokni, Paes Leme STOC18]

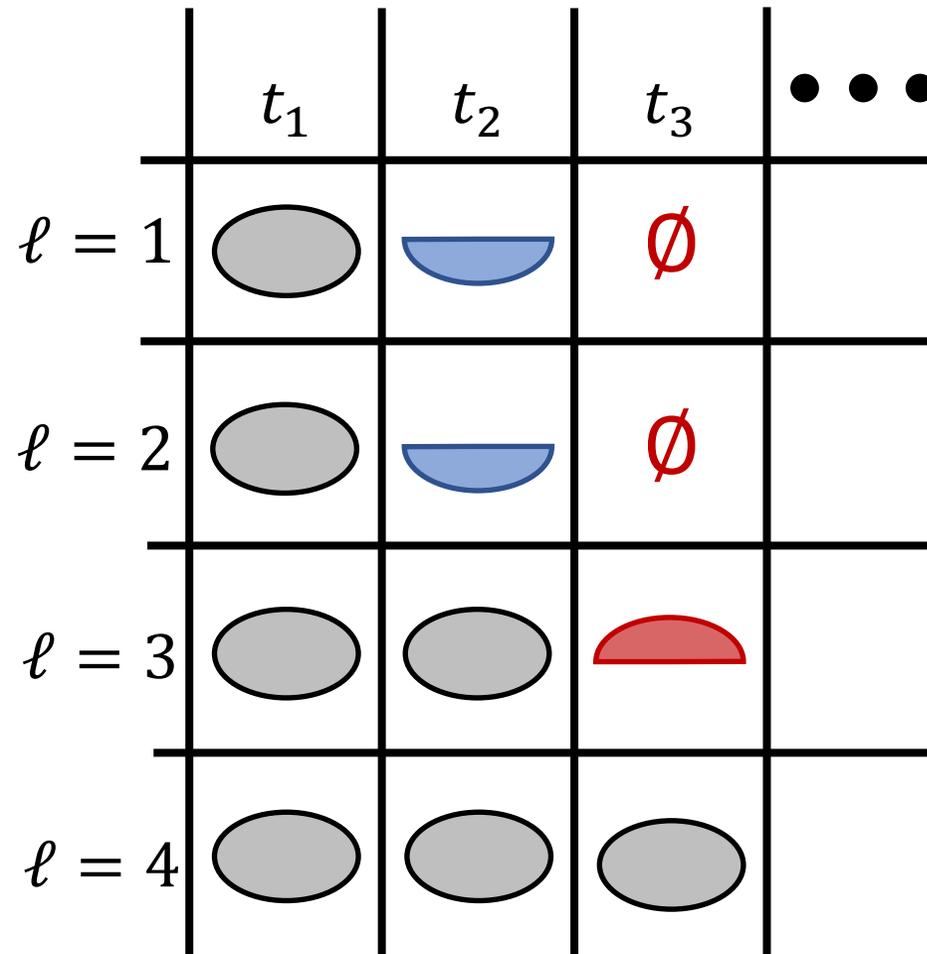
Adaptation to continuous action spaces: [Krishnamurthy, Lykouris, [Podimata](#), Schapire, STOC21]

## Subsampling

- $\log T$  layers
- Layer  $\ell$  corresponds to corruption level:  $C = 2^\ell$
- Sample a layer  $\ell$  to play, with probability  $2^{-\ell}$

## Global Eliminations

- Maintain consistent knowledge sets
- Below layers: more robust  $\rightarrow$  they should inform above layers.



# Incentive-Compatible and Incentive-Aware Learning

! Need to draw both from literature in ML and Mechanism Design.

## Questions to keep in mind for modeling incentive-compatible and incentive-aware learning settings

1. What is the <sup>Decide who to admit in college</sup> goal of the learner, and what the <sup>“Pass” the cutoff and get admitted</sup> goals of the agents?
2. <sup>Candidates' submitted SAT scores</sup> What is observed by the learner regarding the agents?
3. <sup>Times that they can take SAT, amount of money spent for tutoring</sup> What is the agents' ability to respond to decisions that are made?

Stackelberg  
Security Games

Strategic  
Classification

Dynamic Pricing

# The Broader Picture

We have merely scratched the surface!

This Course

## Incentive-Compatible and Incentive-Aware Learning

Examples of topics not covered here: Learning in auctions, learning in games, learning from revealed preferences, learning for matching markets, recommendation systems and many more...

If interested in learning more about Incentive-Compatible Learning:

[Chen, [Podimata](#), Procaccia, Shah, EC18]: algorithms for strategyproof high dimensional linear regression

[Feng, [Podimata](#), Syrgkanis, EC18]: bandit algorithms for learning bidders in online auctions

[Freeman, Pennock, [Podimata](#), Vaughan, ICML20]: strategyproof online prediction algorithms with optimal regret guarantees



Turing



Nala



Jacob



Terra

# Thank You!